# From PNS to TARGET2: the cost of FIFO in RTGS payment systems



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### **Presentation Layout**

- I. Offsetting algorithms
- II. Bilateral optimization
  - a) Standard offsetting algorithms (PNS/TARGET2)
  - b) Non-Fifo advanced algorithms
- III. Multilateral optimization
  - a) Standard offsetting algorithms (PNS/TARGET2)
  - b) Non-Fifo advanced algorithms
- IV. Application 1 : Settlement of an AS after a disaster
- V. Application 2 : Mitigating the consequences of a technical default in the French LVPS PNS

### **Offsetting algorithms**

- Introduced in RTGS to lower the liquidity needs of the system in normal operation
- Can also be used in order to mitigate crisis situation
- Balance computational time / efficiency
- FIFO ("First In First Out") or not FIFO ?
  - PNS and T2 Already not FIFO in many ways
    - MUST in PNS / Normal payments in T2
    - FIFO only valid on a bilateral basis
  - FIFO not algorithm dependant

## Bilateral optimization Standard offsetting algorithms

- Same in PNS and in TARGET2
- 2 equivalent constraints (position and bilateral limit)
- Pure FIFO
- Starts with all payments selected and unselect the last payment of the participant in deficit

	А	В
Position	10	10
Queued payments	500	20
between A and B	20	20
	20	20
	20	20

 No payment settled with PNS/T2 algorithm

- Greedy algorithm (Güntzer et al., 1998)
- Non-FIFO
- Starts with all payments selected
- Unselect all payments from the participant in deficit.
- Reselect the payments from the biggest to the smallest.

	А	В
Position	10	10
Queued payments	500	20
between A and B	20	20
	20	20
	20	20

Success

- Greedy algorithm gives the best solution (in value) for superincreasing payment values.
- If the payment queued are not superincreasing, Greedy will not necessarily provide a good solution.

	Α	В
Position	5	5
Queued navments	140	20
between A and B		20
		30
		100

 No payment settled with Greedy algorithm

- 1st tentative to improve on Greedy:
- Greedy++ : after each Greedy iteration, call a subfunction looking for the best solution using the 10 payments closest to the error. (Test every possibility, 1024 in total)

	А	В
Position	5	5
Queued payments	140	20
between A and B		20
		30
		100

 Greedy++ : error 30, send the 10 closest payments to the subfunction.

Payment 30 is unselected. Success

- 2nd tentative to improve on Greedy:
- Las Vegas Greedy : As in Greedy, consider payments in decreasing order.
- When appropriate, select a payment with a given probability. Run the algorithm several times.

	А	В
Position	5	5
Queued payments	140	20
between A and B		20
		30
		100

Las Vegas Greedy: probability to select 30: 30/(20+20) 25% chance to give the good solution. After 10 tries: 95%

## Bilateral optimization Efficiency of the algorithms: Value

Efficiency test: 2 participants A and B.

- 100 payments queued from A to B. 100 payments queued from B to A.
- Payments generated randomly (lognormal PNS). Average over 5000 distributions.
- Net receiver has no liquidity. Net emitter has α% of the value needed to settle all payments. Measure the settled value as a % of the maximum possible.
- Las Vegas: stop after 5 unsucessful tries in a row.



## Bilateral optimization Efficiency of the algorithms: Volume

Efficiency test: 2 participants A and B.

- 100 payments queued from A to B. 100 payments queued from B to A.
- Payments generated randomly (lognormal PNS). Average over 5000 distributions.
- Net receiver has no liquidity. Net emitter has α% of the value needed to settle all payments. Measure the settled volume as a % of all payments.
- Las Vegas: stop after 5 unsucessful tries in a row.



## Multilateral optimization Standard offsetting algorithms

#### **PNS/T2** multilateral algorithm

- Unselect all payments violating bilateral limits
- Find the participant with the smallest negative virtual position
  - If it exists, inactivate the smallest payment whose
- value is higher than the deficit.
  Otherwise inactivate the payment
- Otherwise inactivate the payment with the highest value.
- Inactivate the most recent payment from this participant.

## Multilateral optimization Standard offsetting algorithms



PNS: The payment of value 100 (the smallest whose value is higher than the deficit 70) is unselected. Failure
 TARGET2: The most recent payment is unselected.

## Multilateral optimization Advanced non-FIFO offsetting algorithms

- Some ideas to improve on the current PNS algorithm
- Favour liquidity transfers towards the center by removing peripherical participants
- Multilateral PNS Las Vegas
  - Same in construction as PNS
  - Starting participant in deficit chosen randomly
  - Payment inactivated chosen randomly (using appropriate probabilities)
- Multilateral Greedy Las Vegas
  - Starting participant in deficit chosen randomly
  - All payments from the considered participant unselected and considered for re-selection in the decreasing order of their value (using appropriate probabilities)

## Multilateral optimization Advanced non-FIFO offsetting algorithms

#### OPM 10-10

- Choose a bank with a negative virtual position
- For each of the outgoing payments of this bank calculate a coefficient depending on:
  - How close the payment is to the deficit of the bank
  - Whether unselecting the payment makes the bank's virtual position positive
  - Whether unselecting the payment creates or amplifies another bank's deficit
- The payment with the highest coefficient is unselected

### Settlement of an urgent AS during a crisis. Scenario

- Settlement of AS after an operational problem.
- Zero liquidity in the system (pre-disaster positions not accessible, f.ex Regional disaster in T2).
- 1 highly urgent ancillary system ("all or nothing", 10 participants) waiting to be settled...
- A certain number of lower importance payments between the AS participants (generated randomly according to a lognormal law m=4.4, σ=1.6)
- Aim: reduce the liquidity needs of the AS participants in a net debit position.

## Settlement of an urgent AS during a crisis. Scenario



- 9 participants with an AS position of +11M
- 1 participant with an AS position of -100M
- Let the multilateral optimisation algorithm select some of the N queued payments in order to reduce the liquidity needs of the participant in a net debit position
- Here liquidity ratio=38%

### Settlement of an urgent AS during a crisis. Results

Liquidity ratio vs algorithm and number of low priority payments available Average over 100 randomly generated distributions



Number of low priority queued payments available

**PNS** structure

#### PNS

- Privately owned large value payment system
- March 2006:
  - 17 participants
  - Between €45 and €90 billions settled per day
  - Around 20 000 payments per day
- Real-time gross system with bilateral limits and optimization algorithms

#### Principle

- The biggest participant faces operational problems
  - It is unable to send payments...
  - But it still receives payments from the other participants
  - And thus turns into a "liquidity trap"
- Simulations
  - Real data used
  - BdF's simulator reproduces exactly the behavior of PNS (data, entry mechanism, optimization algorithms...)
- Measured consequences of the technical default...
  - Increase in settlement delay
  - Rejected payments at the end of the day

#### Advanced algorithms in the case of a default

- Simulations were made to assess the impact of advanced offsetting algorithms on the system, in case of the technical default of the biggest participant (17 march 2006)
  - Normal PNS
  - PNS' FIFO bilateral optimization replaced by Greedy
  - PNS' FIFO bilateral optimization replaced by Greedy++
  - PNS' FIFO bilateral optimization replaced by Las Vegas Greedy
  - PNS' multilateral optimization replaced by M-Greedy-LV
  - PNS' multilateral optimization replaced by M-PNS-LV
  - PNS' multilateral optimization replaced by OPM 10-10

- Results
- 17 March, technical default of the biggest participant
- Various algorithms in replacement of PNS'
  - Bilateral (Las Vegas: stop after 5 unsucessful tries in a row)
  - Multilateral (Las Vegas: stop after 5 unsucessful tries in a row)

#### Rejected payments (value)

#### Rejected payments (volume)



- Results
- 17 March, technical default of the biggest participant
- Various algorithms in replacement of PNS'
  - Bilateral (Las Vegas: stop after 5 unsucessful tries in a row)
  - Multilateral (Las Vegas: stop after 5 unsucessful tries in a row)



#### **Delay indicator (value)**

#### **Delay indicator (volume)**

#### Insight providing example

- Technical default of the biggest participant 17/03/2006
- Focus on the bilateral relation between 2 participants

•	Payments
	rejected at the
	end of the day
	by PNS

- Payments settled with Greedy++
- Payments settled with Greedy

	Α	В
Cash position €	3.5 M	22.5 M
Queued payments between A and B	160 M	1 000 M
	313 M	3 500 M
	956 M	87 M
	1 500 M	
	2 000 M	
	51 M	-
	180 M	and a set

#### Conclusion

- Optimization algorithms can in some cases:
  - Lower the settlement delay
  - Lower the value of the rejected payments
- Two different approaches
  - Non-Fifo algorithms in replacement
  - Non-Fifo algorithms only as a final optimization before rejection
- However this effect is very case-dependent
  - Sometimes the final optimization brings nothing, as PNS algorithms have already settled many payments (2 cases out of 6)

### **General conclusion**

#### Today: liquidity rich systems

- Free providing of intraday liquidity against collateral
- No need for more advanced algorithms in normal operation (although increase in settlement speed non negligible)
- Under special circumstances
  - Technical default / Liquidity crisis
  - Can help mitigate the consequences
    - Settle a highly urgent AS faster
    - Lower the number of rejected payments during a crisis
- Calculation time
  - Greedy faster than PNS, OPM 10-10 not significantly slower, LV 3 times slower
- Non-Fifo advanced algorithms
  - Could be useful in some circumstances
  - Still a lot of room for improvements

### References

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